

Concepts of Experiments at Future Colliders II

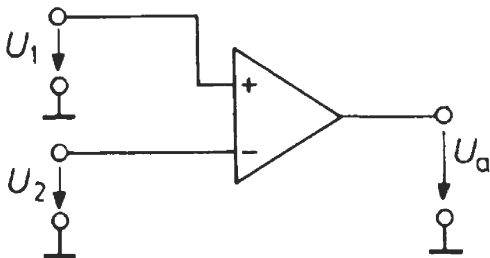
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17.05.2024

From analog to digital signals

Operational amplifiers as comparators

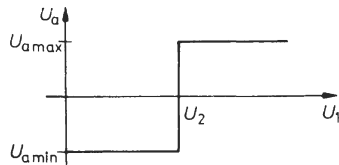
- An operational amplifier saturates when $|U_P - U_N|$ exceeds a small range of values.
- **Comparators** are operational amplifiers where this range has been chosen very small.



In the ideal case:

$$U_a = \begin{cases} U_{a,\max} & \text{for } U_1 > U_2, \\ U_{a,\min} & \text{for } U_1 < U_2. \end{cases}$$

Characteristic curve:



Analog-to-Digital Converter (ADC)

Two basic types of analog-to-digital converters are distinguished.

- Charge-sensing analog-to-digital converter
Measurement of

$$Q := \int_{t_0}^{t_0 + \Delta t} I(t) dt$$

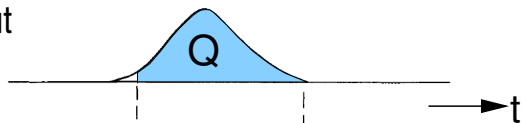
and conversion of the measured value into an integer.

- Amplitude sensing analog-to-digital converter
Measurement of the peak value of a signal $U(t)$ in the interval $[t_0, t_0 + \Delta t]$ and conversion of the measured value into an integer.

Recapitulation of the previous lecture

Wilkinson's method for charge measurement

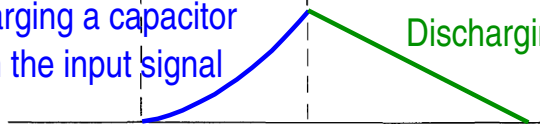
Input



Logical signal defining the time window

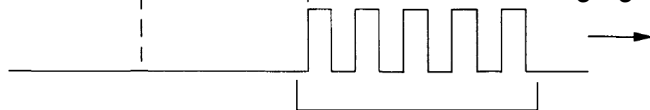


Charging a capacitor with the input signal



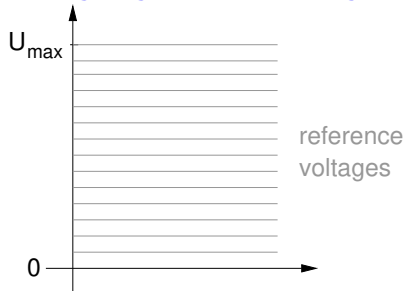
Discharging at constant current

Oscillator for the measurement of the discharging time (prop. to Q)



Number of oscillations = direct measure of the discharging time

Weighing method for signal amplitude measurement



Division of the dynamic range of the analog-to-digital converter into a series of comparison voltages.

Conversion of the results of the voltage comparisons into a bit pattern.

Time-to-Digital Converter (TDC)

Analog signal \rightarrow Comparator \rightarrow Logic signal \rightarrow Time measurement

Simplest approach to time measurement

- Clock generator with a period T smaller than the desired time measurement accuracy.
- Continuous counting of clock cycles. Use a counter with n bits such that $2^n \cdot T >$ (time intervals to be measured).
- Record at which clock cycles n_{Start} and n_{Stop} the start and stop signals have arrived.

$t_{Start} - t_{Stop}$ is then measured as $n_{Start} - n_{Stop}$.

If the counter overflows, one must use $n_{Start} - n_{Stop} + 1$.

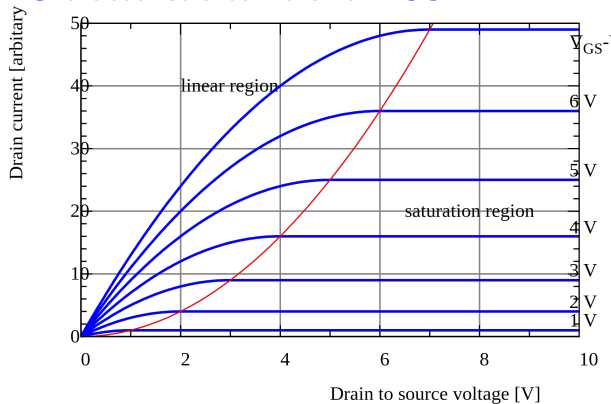
Components for processing digital/logical signals

Logic families

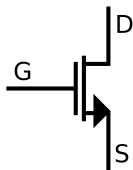
- As mentioned earlier there are different definitions of logic signal levels related to different so-called “logic families”.
- Still in use today (or “popular”):
 - Transistor-transistor logic (TTL) using bipolar transistors.
 - Emitter coupled logic (ECL) using bipolar transistors.
 - Complementary metal oxide semiconductor logic (CMOS) using MOSFETs.

Recapitulation of the previous lecture

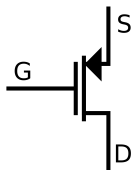
Characteristic curve of a MOSFET



NMOS



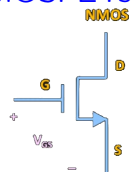
PMOS



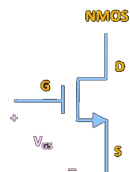
- MOSFETs are operated in saturation mode for logic gates.

Recapitulation of the previous lecture

MOSFETs as switches



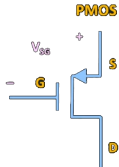
$$V_{GS} < V_T$$



$$V_{GS} > V_T$$



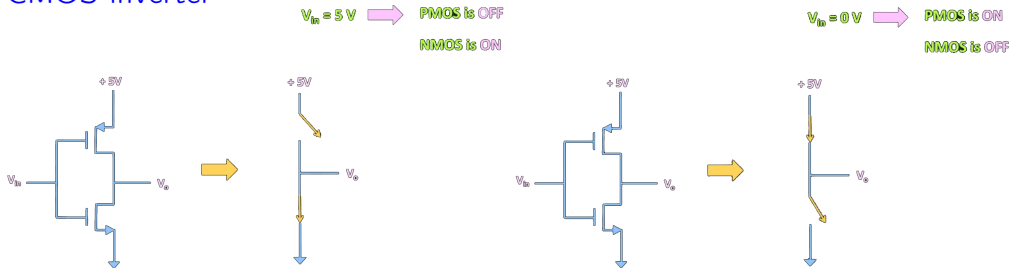
$$V_{SG} < |V_T|$$



$$V_{SG} > |V_T|$$

Recapitulation of the previous lecture

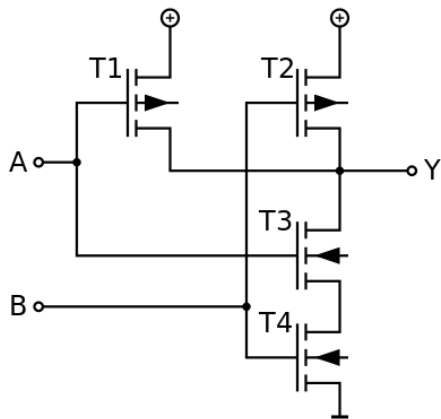
CMOS inverter



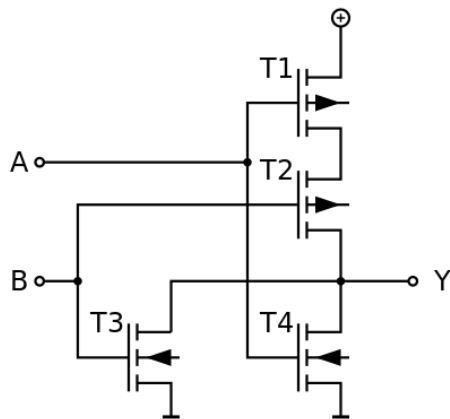
Recapitulation of the previous lecture

CMOS NAND and NOR

$$Y = \overline{A \text{ AND } B}$$



$$Y = \overline{A \text{ OR } B}$$



Logical basic functions

Two States: logical 0 and logical 1.

Logical Basic Functions

- Conjunction: $y = x_1 \wedge x_2 = x_1 \cdot x_2 = x_1 x_2$.
- Disjunction: $y = x_1 \vee x_2 = x_1 + x_2$.
- Negation: $y = \bar{x}$.

Recapitulation of the previous lecture

Rules of Calculation

Kommutatives Gesetz:

$$x_1 x_2 = x_2 x_1$$

Assoziatives Gesetz:

$$x_1(x_2 x_3) = (x_1 x_2)x_3$$

Distributives Gesetz:

$$x_1(x_2 + x_3) = x_1 x_2 + x_1 x_3$$

Absorptionsgesetz:

$$x_1(x_1 + x_2) = x_1$$

Tautologie:

$$x x = x$$

Gesetz für die Negation

$$x \bar{x} = 0$$

Doppelte Negation:

$$\overline{(\bar{x})} = x$$

De Morgans Gesetz:

$$\overline{x_1 x_2} = \bar{x}_1 + \bar{x}_2$$

Operationen mit 0 und 1:

$$x \cdot 1 = x$$

$$x \cdot 0 = 0$$

$$\bar{x} \cdot 1 = \bar{x}$$

$$\bar{x} \cdot 0 = 0$$

$$x_1 + x_2 = x_2 + x_1$$

$$\begin{aligned} x_1 + (x_2 + x_3) \\ = (x_1 + x_2) + x_3 \end{aligned}$$

$$\begin{aligned} x_1 + x_2 x_3 \\ = (x_1 + x_2)(x_1 + x_3) \end{aligned}$$

$$x_1 + x_1 x_2 = x_1$$

$$x + x = x$$

$$x + \bar{x} = 1$$

$$\overline{\overline{x_1 + x_2}} = \bar{x}_1 \bar{x}_2$$

$$x + 0 = x$$

$$x + 1 = 1$$

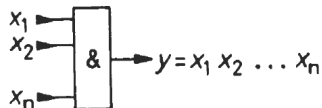
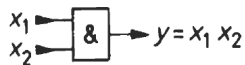
$$\bar{x} + 1 = 1$$

$$\bar{x} + 0 = \bar{x}$$

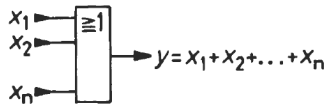
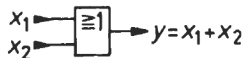
Recapitulation of the previous lecture

Switching elements for logical basic functions

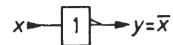
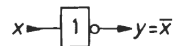
Conjunction AND Gate



Disjunction OR Gate



Negation NOT Gate



Method of disjunctive normal form

To establish more complex logical functions, one can use the so-called **disjunctive normal form**.

n input variables x_1, \dots, x_n . 1 output variable y .

1. Set up a table listing all possible input values along with the desired output value. This table is also called a **truth table**.
2. Identify all rows in the truth table where $y = 1$.
3. For each of these rows, form the conjunction of all input variables; for $x_k = 1$, substitute x_k , otherwise \bar{x}_k .
4. The sought function is obtained by forming the disjunction of all found product terms.

Example of exclusive OR

Truth Table

Row	x_1	x_2	y	
1	1	1	0	
2	1	0	1	$\rightarrow x_1 \cdot \bar{x}_2 =: K_2$
3	0	1	1	$\rightarrow \bar{x}_1 \cdot x_2 =: K_3$
4	0	0	0	

Result

$$y = K_2 + K_3 = (x_1 \cdot \bar{x}_2) + (\bar{x}_1 \cdot x_2).$$

Example of a 1-of-4 decoder

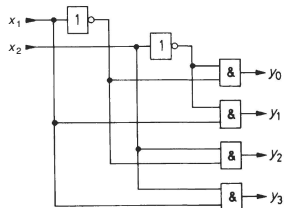
Truth table

Row	x_1	x_2	y_3	y_2	y_1	y_0
1	0	0	0	0	0	1
2	0	1	0	0	1	0
3	1	0	0	1	0	0
4	1	1	1	0	0	0

Result

$$y_0 = \bar{x}_1 \cdot \bar{x}_2. \quad y_1 = \bar{x}_1 \cdot x_2. \quad y_2 = x_1 \cdot \bar{x}_2. \quad y_3 = x_1 \cdot x_2.$$

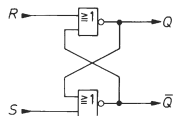
Circuit



Derived Basic Functions

$$x_1 \text{ NOR } x_2 := \overline{x_1 + x_2} = \bar{x}_1 \cdot \bar{x}_2. \quad \text{[NOR symbol]} \rightarrow Y$$
$$x_1 \text{ NAND } x_2 := \overline{x_1 \cdot x_2} = \bar{x}_1 + \bar{x}_2. \quad \text{[NAND symbol]} \rightarrow Y$$

Flip-Flop



$$Q = \overline{\bar{Q} + R}.$$

$$\bar{Q} = \overline{S + Q}.$$

S	R	Q	\bar{Q}
0	0	Q_{-1}	\bar{Q}_{-1}
0	1	0	1
1	0	1	0
1	1	(0)	(0)

(bisheriger Zustand)
(Zurücksetzen: 0|1)
(Setzen: 1|0)

Setting $S = R = 1$ results in $Q = \overline{\bar{Q} + 1} = \bar{1} = 0$ and $\bar{Q} = \overline{1 + Q} = \bar{1} = 0$.
Subsequently setting $R = 0$ and $S = 0$ simultaneously leaves the output state undefined.

$$Q = \overline{\bar{Q} + 0} = \bar{\bar{Q}} \text{ can be 0 or 1.}$$

$$\bar{Q} = \overline{Q + 0} = \bar{Q} \text{ can be 0 or 1.}$$

$\Rightarrow R = S = 1$ is generally prohibited.

Fundamentals of statistical treatment of experimental data

Introductory example: beam energy measurement

Example: Measurement of the energy of a monoenergetic particle beam.

Notations

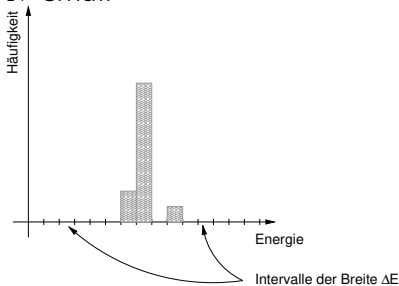
E_S : actual beam energy.

N : number of measurements of beam energy.

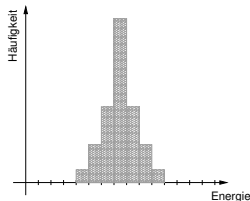
E_k : result of the k -th measurement of beam energy.

Frequency Distributions

N small

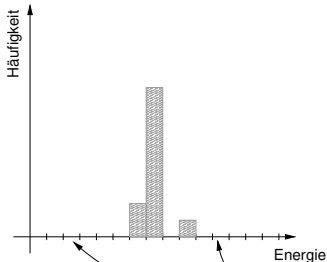


N large

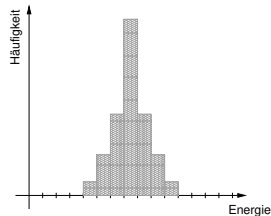


Introductory example: beam energy measurement

N small



N large

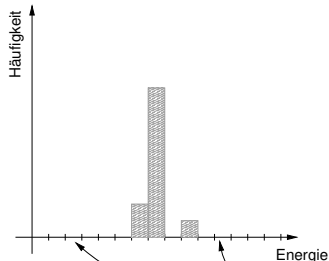


Intervalle der Breite ΔE

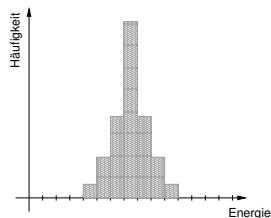
- When N is large, repeating the N measurements yields (nearly) the same frequency distribution.
- In the limit $N \rightarrow \infty$, the frequency distribution converges to the probability distribution for the outcome of the measurement.

Introductory example: beam energy measurement

N small



N large



Intervalle der Breite ΔE

- The probability of measuring E_k when the beam energy is E_S depends on the value of E_S and the measurement method. If one knows the probability function $p(E_k; E_S)$, one can determine E_S from the measurement of the frequency distribution.
- In practice, $p(E_k; E_S)$ is only partially known, and one tries to infer $p(E_k; E_S)$ from the measured frequency distribution, which provides an estimate of E_S . In statistics, methods are employed to infer the underlying probability distributions from frequency distributions.

- A physical measurement is a **random process**.
- A measured quantity x , which represents the outcome of a random process, is called a **random variable** or **random quantity**.
- Any function of x is also a random variable.
- If the random variable can only take discrete values, there is a probability for the occurrence of each of these values, which is the **probability function**.
- For random variables with continuous range of values, the **probability density** $p(x)$ replaces the probability function. Let Ω be a measurable set of possible values of x , whose measure is greater than zero. Then

$$\int_{\Omega} p(x) dx$$

is the probability of observing a value $x \in \Omega$.

The mathematical field of probability theory is based on [Kolmogorov's Axioms](#).

Kolmogorov's Axioms

Let Σ denote a set of events.

1. For every event $A \in \Sigma$, the probability of the occurrence of A is a real number $p(A) \in [0, 1]$.
2. The certain event $S \in \Sigma$ has probability $p(S) = 1$.
3. The probability of the union of countably many incompatible events is equal to the sum of the probabilities of the individual events. Here, events A_k are [incompatible](#) if they are pairwise disjoint, i.e., $A_k \cap A_\ell = \emptyset$ for all $k \neq \ell$.

Characteristics of probability distributions

Remark. In this section, we consider probability densities. Probability functions of discrete variables are also covered if one considers the δ -distribution as a probability density.

Nomenclature. D : Range of values of a random variable $x = (x_1, \dots, x_n)$.
 $p(x)$: Probability density of x .
(D is the domain of p .)

Definitions

The **expectation value** of x , $E(x)$ (also $\langle x \rangle$), is defined as

$$E(x) := \int_D x \cdot p(x) dx.$$

The **covariance matrix** $cov(x_k, x_l)$ is defined as

$$cov(x_k, x_l) := \langle (x_k - \langle x_k \rangle) \cdot (x_l - \langle x_l \rangle) \rangle .$$

The diagonal element $cov(x_k, x_k)$ is called the **variance of x_k** , $Var(x_k)$, and $\sqrt{Var(x_k)}$ is the **standard deviation** $\sigma(x_k)$.

Expectation value of a function of a random variable

- A function $f(x)$ is also a random variable.

$$\langle f \rangle = \int_D f(x)p(x)dx.$$

- If $f(x) = f(x - \langle x \rangle + \langle x \rangle)$ is significantly different from 0 only for small values of $|x - \langle x \rangle|$, one can approximate $f(x)$ by

$$f(\langle x \rangle) + \left. \frac{df}{dx} \right|_{\langle x \rangle} \cdot (x - \langle x \rangle)$$

Then

$$\begin{aligned} \langle f \rangle &\approx \left\langle f(\langle x \rangle) + \left. \frac{df}{dx} \right|_{\langle x \rangle} \cdot (x - \langle x \rangle) \right\rangle \\ &= \langle f(x) \rangle + \left\langle \left. \frac{df}{dx} \right|_{\langle x \rangle} \cdot (x - \langle x \rangle) \right\rangle \\ &= f(\langle x \rangle) + \left. \frac{df}{dx} \right|_{\langle x \rangle} \cdot (\langle x \rangle - \langle x \rangle) = f(\langle x \rangle). \end{aligned}$$

Variance of a function of a random variable

Special Case: $f(x) \in \mathbb{R}$.

$$\begin{aligned} \text{Var}(f) &= \langle (f - \langle f \rangle)^2 \rangle = \langle [f - f(\langle x \rangle)] \rangle \\ &\approx \left\langle \left[\sum_{k=1}^n \frac{df}{dx_k} \Big|_{\langle x \rangle} \cdot (x_k - \langle x_k \rangle) \right]^2 \right\rangle \\ &= \left\langle \left[\sum_{k,\ell=1}^n \frac{df}{dx_k} \Big|_{\langle x \rangle} \frac{df}{dx_\ell} \Big|_{\langle x \rangle} \cdot (x_k - \langle x_k \rangle) \cdot (x_\ell - \langle x_\ell \rangle) \right] \right\rangle \\ &= \sum_{k,\ell=1}^n \frac{df}{dx_k} \Big|_{\langle x \rangle} \frac{df}{dx_\ell} \Big|_{\langle x \rangle} \cdot \langle (x_k - \langle x_k \rangle) \cdot (x_\ell - \langle x_\ell \rangle) \rangle \\ &= \sum_{k,\ell=1}^n \frac{df}{dx_k} \Big|_{\langle x \rangle} \frac{df}{dx_\ell} \Big|_{\langle x \rangle} \cdot \text{cov}(x_k, x_\ell), \end{aligned}$$

which is the well-known [error propagation formula](#).

Examples of important probability distributions

The binomial distribution

- The **binomial distribution** gives the probability of observing n_k events out of a total of N events when ν_k events are expected:

$$p(n_k; \nu_k) = \binom{N}{n_k} \left(\frac{\nu_k}{N}\right)^{n_k} \left(1 - \frac{\nu_k}{N}\right)^{N-n_k}.$$

- With $p := \frac{\nu_k}{N}$, one obtains from

$$\begin{aligned} 0 &= \frac{d}{dp} 1 = \frac{d}{dp} \sum_{n_k=0}^N \binom{N}{n_k} p^{n_k} (1-p)^{N-n_k} \\ &= \sum_{n_k=0}^N \binom{N}{n_k} [n_k p^{n_k-1} (1-p)^{N-n_k} - (N-n_k) p^{n_k} (1-p)^{N-n_k-1}] \\ &= \frac{1}{p} \langle n_k \rangle - \frac{1}{1-p} \langle N - n_k \rangle = \left(\frac{1}{p} + \frac{1}{1-p}\right) \langle n_k \rangle + \frac{N}{1-p} \\ &= \frac{1}{p(1-p)} \langle n_k \rangle + \frac{N}{1-p} \Leftrightarrow \langle n_k \rangle = N \cdot p = N \cdot \frac{\nu_k}{N} = \nu_k. \end{aligned}$$

- Using the same calculation trick, one obtains $\text{Var}(n_k) = \nu_k(1 - \frac{\nu_k}{N})$.